# CP/ $\mathrm{B}_{2}$ DUnAnic DeBuccinc USERS MANUAL 



VETROZ

CP/M DYNAMIC DEBUGGING TOOL (DDT) USER'S GUIDE

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User ${ }^{\circ} \mathrm{s}$ Guide
I. Introduction.

The DDI trogram allows dynamic interactive testing and debugging of programs generated in the $C P / M$ enviromment. The debugger is initiated by typing one of the following commands at tie CP/M Console Command level

```
DDI
DDI filename.HEX
DDI filename.COM
```

where "filename" is the name of the program to be loaded and tested. In both cases, the DDI program is brought into main memory in the place of the Console Command Processor (refer to the CP/M Interface Guide for standard memory organization), and thus resides directly below the Basic Disk Operating System portion of CP/M. The BDOS starting address, which is located in the address field of the JMP instruction at location 5 H, is altered to reflect the reduced Transient Program Area size.

The second and third foms of the DDI command shown aiove perform the same actions as the first, except there is a subsequent automatic load of the specified iEX or COM file. The action is icientical to the sequence of commands

```
DDI
Ifilename.HEX or Ifilename.COM
R
```

where the $I$ and $R$ commands set $u$ pond read the specified program to test (see the explanation of the $I$ and $R$ commands below for exact details).

Upon initiation, DDI prints a sign-on message in the format
DDI VER m.m

Following the sign on message, DDT prompts the operator with the cinaracter "o" and waits for incut comands from the console. The operator can trye any of several single character commands, terminated by a carriage return to execute the comand. Each line oń input can be linemedited using the standard C8/M conerols

| subout | senove the last character typed |
| :---: | :---: |
| ctlou | semove the entire line, ready for retypi |
| ctloc | Systern seboot |

Any command can be up to 32 characters in lenath (an automatic carriage return is inserted as tise 3350 caractay), where the first coaracter determines the comand trpe

|  | enter assemoly lanquage memonics with operancis |
| :---: | :---: |
| D | display memory in hexadecimal and ASCII |
| E | sill memory with constant data |
| G | begin execution with optional breakpoints |
| I | set up a standard input file controi block |
| $\underline{L}$ | list memory usirg assembler menonics |
| M | move a memory segment fram scurce to destination |
| 8 | gead grogram for subseguent testing |
| 5 | substitute memory values |
| T | srace program execution |
| \% | mbraced prooram monitoring |
| 8 | examine and cotionally alter the cry state |

The camand ciaracter, in some cases, is followed by zero, one, two, or three hexadecimal values which are separated by commas or single blank cinaracters. All DCT numeric output is in hexadecimal form. In all cases, the commancis are not execused until the carsiage securn is typed at the end of the comand.

At any point in the debug fun, the operator can stop execution of DDT using either a ctl-G or GO (jme to location 0000H), and save the current memory imace using a SAVE command of the form

SAve n filename.COM
where $n$ is tiee numier or pages (256 byte blocks) to be saved on disk. The number of blocks can be cetermined by taking the hian orcier byte of tine top load adaress and converting this number to cecimal. For example, if tre highest address in the Transient Prooram Area is 1234 H then the numieer of pages is 12 H , or 18 in cecimal. Thus the cperator could troe a cel-c during the cebug run, returning to the Console Processor level, followed by

SAVE 18 K.COM
The menofy imace is saved as X.COM on the diskettミ, and can be direcrly executed by simply troing tine name $X$. If further testing is required, the memory imace can be recalled by tyoing
which reloads previously saved program from loaction 100H through page 18 (12FFH). The machine state is not a part of the COM file, and thus the program must be restarted fram the beginning in order to properly test it.
II. DDI COMMANDS.

The individual commands are given below in some detail. In each case, the operator must wait for the prompt character ( - ) before entering the command. If control is passed to a program under test, and the program has not reached a breakpoint, control can be returned to DDT by executing a FST 7 from the front panel (note that the rubout key should be used instead if the program is executing a $T$ or $U$ command). In the explanation of each command, the command letter is shown in some cases with numioers separated by cammas, where the numbers are represented by lower case letters. These numbers are always assumed to be in a hexadecimal radix, and from one to four digits in length (longer rumbers will be automatically truncated on the right).

Many of the commands operate upon a "CPU state" which corresponds to the program uncer test. The CPU state holds the registers of the program being debugged, and initially contains zeroes for all registers and flags except for the pogram counter (P) and stack pointer (S), winich default.to 100H. The program counter is swseguently set to the starting address given in the last record of a HEX file if a file of this form is loaded (see the $I$ and $R$ commands) .

1. The A (Assemble) Command. DDI allows inline assembly language to be inserted into the current memory image using the A command which takes the form

## As

where $s$ is the hexadecimal starting address for the inline assembly. DDT prompts the console with the address of the next instruction to fill, and reads the console, looking for assembly language memonics (see the Intel 8080 Assembly Language Reference Card for a list of anemonics), followed by register references and operands in absolute hexadecimal form. Each sucessive load address is printed before reading the console. The A command terminates when the first empry line is irput from the console.

Upon completion of assembly language input, the operator can review the memory segment using the DDT disassembler (see the I command).

Note that the assembler/disassembler portion of DDT can be overlayed by the transient program being tested, in which case the DDT program responcis with an error condition when the $A$ and $L$ commands are used (refer to Section IV)。
2. The D (Display) Command. The D command allows the operator to view the contents of memory in hexadecimal and ASCII fommes. The forms are

$$
\begin{aligned}
& D \\
& D S \\
& D S, E
\end{aligned}
$$

In the first case, memory is displayed from the current display address (initially 100H), and continues for 16 display lines. Eaci display line takes the fom shown below
aaaa ob b b b b b bb bo boccccaccaccccc
where aaaa is the display adaress in hexadecimal, and bo represents dara present in memory starting at aaaa. The ASCII characters starting at aaaa are given to the right (represented by the sequence of $c^{\circ}$ s), where non-graciic characters are trinted as a geriod (o) symbol. Note that both upper and lower case alphabetics are displayed, and thus will appear as upper case symbols on a console device that sucpores only upper case. Each display line gives the values of 16 bytes of data, except that the first line displayed is truncated so that the next line begins at an address which is a multiple of 16.

The second form of the $D$ camand shown above is similar to the fisst, excopt that the display address is first set to address $s$. The third form causes the display to continue from address $s$ tirough address E. In all cases. the display addeess is set to the first adoress not displayed in this cammand, so that a continuing display can be accomplisined by issuirg successive $D$ camands with no explicit addresses.

Excessively long displays can be aborted by gushing the rubout key.

3o. The $E$ (Fili) Command. The $E$ command takes the fomm

$$
E s, F, C
$$

winere $s$ is the starting acdress, $f$ is the final address, and $c$ is a hexadecimal byte constant. The effect is as follows: DDT stores the constant $c$ at address $s$, increments the value of $s$ and tests against $f$. If $s$ exceeds $f$ then the operation terminates, otherwise the aceration is repeated. Thus, the fill command can be used to set a memorl block to a specific constanc value.
4. The $G$ (GO) Commanc. Program execution is started using the $G$ comand, with up to toso cotional breakpoint adcresses. The $G$ command takes one ot the EOms

$$
\begin{aligned}
& G \\
& G: b
\end{aligned}
$$

$$
\begin{aligned}
& G, b, c \\
& G, b \\
& G, b, c
\end{aligned}
$$

The first form starts execution of the program under test at the current value of the program counter in the current machine state, with no breakpoints set (the only way to regain control in DDT is through a RST 7 execution). The current program counter can be viewed by typing an $X$ or $X P$ conmand. The second form is similar to the first except that the program counter in the current machine state is set to address $s$ before execution begins. The third form is the same as the second, except that orogram execution stops when address $b$ is encountered ( $b$ must be in the area of the program under test). The instruction at location $b$ is not executed when the breakpoint is encountered. The fourth fom is identical to the third, except that two breakpoints are specified, one at $b$ and the other at $c$. Encountering either breakpoint causes execution to stop, and both breakpoints are subsequently cleared. The last two forms take the program counter from the current machine state, and set one and two breakpoints, respectively.

Execution continues fram the starting address in real-time to the next breakpint. That is, there is no intervention between the starting address and the break address by DDT. Thus, if the program under test does not reach a breakpoint, control cannot return to DDT without executing a RST 7 instruction. Upon encountering a breakpoint, DDT stops execution and types

* ${ }^{\text {C }}$
where $d$ is the stop address. The machine state can be examined at this point using the X (Examine) conmand. The operator must specify breakpoints wich differ from the program counter address at the beginning of the $G$ command. Thus, if the current program counter is 1234 H , then the conmands
G. 1234
and
G400.400
both produce an immediate breakpoint, without executing any instructions whatsoever.

5. The I (Input) Command. The I command allows the operator to insert a file name into the default file control block at 5CH (the file control block created by CP/M for transient programs is placed at this location; see the CP/M Interface Guide). The default FCB can be used by the program under test as if it had been passed by the CP/M Console Processor. Note that this file neme is also used by DDT for reading additional HEX and COM files. The fom of the I command is

> Ifilename

Or

## Ifilename.filletyce

If the second form is used, and the Eiletype is either HEX or COM, then subsequent R commands can be used to read the pure binary or hex format machine cocie (see the $R$ command for further details).
б. The L (List) Command. The $L$ command is used to list assembly language mnemonics in a particular program region. The foms are

$$
\begin{aligned}
& L \\
& \text { Ls } \\
& \text { Ls,E }
\end{aligned}
$$

The first camand lists twelve lines of disassemoled machine code from the current list address. The second form sets the list adoress to $s$, and then lists twelve lines of code. The last fom lists disassented code from s through address f. In all three cases, the list address is set to the next unlisted location in preparation for a subseguent i command. upon encountering an execution breakpoint, the list address is set to the current value of the program counter (see the $G$ and $T$ commands). Aceain, long trpeouts can be aborted using the rubout key during the list grocess.
7. The $M$ (Move) Comurand. The $M$ commard allows block movement of procram or cata areas frcm one location to another in memory. the form is
Ms,f,d
where $s$ is the start address of the move, $f$ is the final address of the move, and $d$ is the destination address. Data is first moved fron $s$ to $d$, and both acicresses are incremented. if $s$ excesds finen the move operation stops, otherwise the move operation is repeated.
8. The $R$ (Read) Command. The $R$ command is used in conjunction with the I command to read CCM and HEX files from tine diskette into the transient orocram area in oreparation for the debug run. The foms are

## R <br> Rb

where $b$ is an cotional bias adciress which is 三died to each orogram or data acciress as it is loaced. The load qeration must not overwrite any of the system garameters $f=0 \mathrm{~m}$ 000H through $0 F F H$ (i.e., the first page of memory). If b is anitted, then $b=: 0000$ is assumed. The R command rectuires a orevicus $I$ commard, spacifying the name or a HEX or COM ilie. The load adcress for each record is obtained from each individual aEX record, while an assumed lead address of 100 H is taken for COM files. Note that any number of R commands can be issued following the I command to rereed the orocram under test,
assuming the tested program does not destroy the default area at 5CH. Further, any file specified with the filetype "COM" is assumed to contain machine code in pure binary form (created with the LOAD or SAVE command), and all others are assumed to contain machine code in Intel hex format (produced, for example, with the ASM command).

Recall that the command
DDT filename.filetype
which initiates the DDT program is equivalent to the commands

```
DDr
-Ifilename.filetype
-R
```

Whenever the $R$ command is issued, DDT responds with either the error indicator "?" (file cannot be cpened, or a necksum error occurred in a HEX file), or with a load message taking the form

NEXT PC
popn pap
where nnnn is the next address following the loaded program, and pppp is the assumed program counter (100H for COM files, or taken from the last record if a HEX file is specified).
9. The $S$ (Set) Command. The $S$ command allows memory locations to be examined and cotionally altered. The form of the command is

## Ss

where $s$ is the hexadecimal starting adaress for examination and alteration of memory. DDT responds with a numeric orompt, givina the memory location, along with the data currently held in the memory location. If the operator types a carriage return, then the data is not altered. If a byte value is typed, then the value is stored at the frompted address. In either case, DDT continues to prampt with successive addresses and values until either a period (0) is troed by the perator, or an invalid incut value is detected.
19. The $T$ (Trace) Command. The $T$ command allows selective tracing of program execution for 1 to 65535 program steps. The forms are

T
In
In the first case, the Cory state is displayed, and the next program step is executec. The pooram teminates immediately, with the termination address

## *hhhh

where hinin is the next address to execute. The display address (used in the D command) is set to the value of $H$ and $L$, and the list acdress (used in the $I$ command) is set to hinh. The CPU state at program termination can then be examined using the X command.

The second form of the $T$ command is similar to the first, except that execution is traced for $n$ steps ( $n$ is a hexadecimal value) before a program breakpoint is occurs. A breakpoint can be forced in the trace mode by trying a rubout character. The CPU state is displayed before each program step is taken in tracs mode. The fomat of tise display is the same as described in the $X$ command.

Note that program tracing is discontinued at the interface to $6 P / M$, and resumes after return from $C P / M$ to the program under test. Thus. $C 2 / M$ functions witich access I/O devices, such as the diskette drive, run in real-time, avoiding I/O timing oroblems. Programs running in trace mode execute approximately 500 times slower than real time since DDF gets control after eaci user instruction is executed. Interrupt grocessing routines can be traced, but it must be noted that cmmards which use the breakpoint facility (G, T. and U) accomplish the break using a RST 7 inseruction, wich means that the tested program cannot use this interrupt location. Further, the trace moce always runs the tested orogram with interrupes enabled, which may cause problems if asyncironous interrupts are received during tracira.

Note also that the operator sinould use the rubout key to get control back to DCT during trace, rather than executing a RST 7, in order to ensure that the trace for the arrent instruction is completed berore intermuption.
11. The $U$ (Untrace) Command. The $U$ command is identical to the $I$ command except that intemediate frogram steps are not displayed. The untrace mode allows from 1 to 65535 (gFFFFH) steps to be executed in monitored moce, and is used grincipally to retain control of an executing proaram while it reaches steady state conditions. All conditions of the $T$ command appiy to the $U$ command.
12. The $X$ (Examine) Command. The $X$ commanc allows selective display and altaration of the current CPU state for the eroçam uncer test. The forms are
x

Xr
where $r$ is one ci the 8080 CFU registers

$$
\begin{array}{lll}
C & \text { Carry Flag } & (\theta / 1) \\
Z & \text { Zero Flag } & (\theta / 1)
\end{array}
$$

| M | Minus Flag | (0/1) |
| :---: | :---: | :---: |
| E | Even Parity Flag | (0/1) |
| I | Interdigit Carry | (0/1) |
| A | Accumulator | ( 0 -FF) |
| 8 | BC register pair | (0-FFFF) |
| D | CE register pair | (D-FFFF) |
| H | HL reaister pair | ( 0 -FFFF) |
| S | Stack Pointer | ( 0 -FFFF) |
| P | Procram Counter | ( 0 -FFFF) |

In the first case, the CPU register state is displayed in the fomat

where $f$ is a or 1 flag value, bo is a byte value, and dadd is a double byte quantity corresponding to the register pair. The "inst" field contains the disassembled instruction which occurs at the location addressed by the CPU state's program counter.

The second form allows display and optional alteration of register values, where $r$ is one of the registers given above ( $C, Z, M, E, I, A, B, D, H, S$, or P). In each case, the flag or register value is first displayed at the console. The DDI program then accepts input from the console. If a carriage retiun is typed, then the flag or register value is not altered. If a value in the groper range is apped, then tine flag or register value is altered. Note that BC, DE, and HL are displayed as reaister pairs. Thus, the operator tipes the entire register pair when $3, C$, or the CC pair is altered.
III. IMPLEMENTATICN NOTES.

The organization of DDT allows certain non-essential portions to be overlayed in order to gain a larger transient program area for debugging large programs. The DDI procram consists of two parts: the DDI nucleus and the assembler/aisassembler module. The DDT nucleus is loaded over the Console Command Processor, and, althouch loaded with the DDT nucleus, the assembler/disassembler is overlayable unless used to assemble or disassemile.

In particular, the BDOS address at location 6ii (adidress field of the JMP instruction at location 5if) is modified by DDT to acicress the base location of the DDT nucleus which, in turn, contains a JMP instruction to the BDOS. Thus, programs winci use this address field to size memory see the logical end of memory at tine base of the DDT nucleus rather than the base of the BCCS.

The assemoler/disassembler module resices directly below the DDT nucleus in the transient rooram area. If the $A, L$, $T$, or $X$ commands are used during the deiucucing process then the DDT program again alters the acciress field at $6 i$ to irclude this module, thus further reducinc tine loaical end of memory. If a program loads beyond the beginning of the assemoler/disassembler module, the $A$ and $L$ comands are lost (their use produces $a$ "?" in response), and the
trace and display (T and $X$ ) commands list the "inst" field of the display in hexadecimal. bather than as a cecoded instruction.

## IV. EN EXAMPRE.

The following example shows an edit, assemble, and cebug for a simple program which reads a set of data values and detemines the largest value in the set. The largest value is taken from the vector, and stored into "LARGE" at the termination of the program



AEM SiAN Start Assembler
EPMA ASSEMBLER－UER 1.0

5122
Gधご リSE FACTOF
END IF ASSEMBLO
Assemblu Cumplete－Look at Pregram Lesting
THFE SCAH．FRNZ
Code identrs


0192 aع日品 MYI

B．LEM ：EMCTH OF भECTON TO SCAN
H，UEGT．SASE DF YEGTOR
ala7 75
$\operatorname{igof}$
Mu＇
EU8
H，M VET リA！UE

0189 D2EDA！
 NEW LABCEET UALUE．ETGEE ：T，TGE
B！日C $4 F$
815D 23 MUY $\quad \therefore A$

31月E 85

DER E MEE IJ SCHAO
日1日F C2要7！
． N 2
LOOF FOE MNOTHEE

0112 33
911．3 322191
9：16 c3Ag80
Cede／date lisin．
Truncated TEET INTA
0119 62954 5305はECT：
日986 $=~ \longleftrightarrow ~ L E N$
－121 Value 㝖
0：22 Equate
GARGE：

EEU EOYECT ：ENGT母 IS $\quad \therefore \quad$ GAFGEST UGGUE DN E\％！T
A :

BIT SCAN．HE：STart Deiougger using hex format machive code
1BK DET पER 1．G
MEXT PE
a：${ }^{3!}$ last load address +1
一资
hext insinuction
 to execute at - 兑 P $\qquad$ Examine vegistors before deloug run

－Laok at vesisters aguin


$-1$| 160 |
| :--- |

Neat instruction

to crecure at $P_{e}=160$
－its，exier inline assembly mocie to clance $H_{k e}$ Jup $\frac{1}{0} 0000$ linto a RST？which
 will Cause the program under est to retam to OOT if 116 ＇r
Q1：73（single carnase retern stops assembe mode）
－1115，List code at 13 to to ceeck that RST 7 wous properig inseremen


-x. $\quad$ Look at resisters

-I: Execute Program for one step. initial CPU state, before 7 is executed
 -I2 Trice one ster again (note os in B) automatic break point
 -I Trace again (Register Cis cleared)
 -I3) Trace three steps


 - I111 2 Display memory starting at 119 H .












- $3 \quad$ Current CPu state ?

-TE Trace $S$ seeps from Current CPU State




 -1153 Trace withasit listing intermediate states



－it Run Program from current PC unill completion（iu real－time） －a1：breakpoint at $116 \mathrm{H} /$ coused be exeaning RST 7 in macinne cade
－i，ceustute at end of proarm

－X：？Examinc and change Prcaram counter
$\bar{F}=4115 \xrightarrow{380}$
－关
C日2 MaEd！
ile
ind 10 （hexadecmal）sleps
－A！is，Insert a inot patch＂，vito
：1：9 Je ：ity，the machure code
to chanar the Juc to Je
－ing Siop SDT so－inat a version of the pateched progrom can be saved

Programshould have moved the value from $A$ ：nio $C$ since $A>C$ ． Since this cade was not executén， it appears that the JNC should have been a JC insiruction



NEET FC

－L！日星 List some Code


时
¢！ 58 5us
边的
$\square$ 182


- I1B Trace to see how parched version operates Date is moved from $A \frac{1}{O} C$
















- 


-15.1is, Run from current PC and breakpoint at 108 H
*5168
一名
-next data item
 $-12$

Single Step for a few cycles
 -i,
 - X

-i, Run to Completion

* 9 ! ! 5
$-\underset{-1}{-1}$

- Sics look at the value of "LARGE"
©121 83, Wrens Value!

```
4:220%%
4135 2%%
5124 2%,
W125 4 4,
Ni:5 位, End of the S Command
\12% %E - % 
-1.6星
```



```
Rovins the Cole
-L
\1!j sia glet
*1l6 RST 9%
g1/7 NOP
0118 NOP
#1!g STA% =
G1:A NOf
G!!8 !NR S
#1!0 INK 5
GIII DCR 3
G:IE M4! g.a!
ब!टQ DCR 3
-iF
F=m1:E !ag; Reset the PC
-Iz Sivale Step, and unted dam values
```



```
-iz
```



```
-I% Sount set laraar'set
```



```
I:
```



EATE 1 SCAN．COM，
MODT SCAN．COM，

Save memory image
Restant DTT
1Е゙K DET VEK 1．
HEXT FE
－ 2940108

$\bar{F}=\sin \mathrm{A}_{2}$


－G．1．16．Run from look to cimpletion
＊诣！！
－足：Lookat Carry（acenertal Iypo）
42
－iz Look at ciul state

－Ei2：？book at＂Large＂－it appears to be correct．
－ 15185
ज1E

－TD 5 Step DDT
EI SGAM．ASH，Re－edit the scurce prearam，and make boith changes
$-\frac{n \operatorname{BiE} z}{2+5 z}$



＊ 2

NFGUNI JUMP PF LARGER YALUE NOT FOUND
MFOUHT ；JMMP IF LAPEEE VALUE VOT GOUNG
－


$$
\begin{aligned}
& \text { GFM AESEMELEE - リER } 1 \text { a MRX To disk } A \\
& \text { - : こ } \\
& \text { onget use =actor } \\
& \text { ENG OF ASEEMELY } \\
& \begin{array}{l}
\text { nex to disk } A \\
\text { print to } \\
\text { (selects us print file) }
\end{array}
\end{aligned}
$$



$$
\begin{aligned}
& \text { AEMTF: }
\end{aligned}
$$

$$
\begin{aligned}
& \text { - L1! }{ }^{2}
\end{aligned}
$$

a： 15 imp ang check ter ensure end is still at 116 H 5l：シís 三
a！：A NEE

－（rubart）
－ising：！！，Go from beqiuning wift breakpoint at and －it：$:$ breakpoint rearleá
－I：E！，LOOk at＂LARSE＂
Correct vaiur compurtei



－（rubouit）aboris long typecut
－it Stop DOT，debua session compleie

)

